

Threads and Scheduling

Simon Peter

Systems Group Dept. Computer Science ETH Zurich - Switzerland

Littere

Eidgenössische Technische Hochschule Zürich Swiss Federal Institute of Technology Zurich





- Threads
 - Concepts
 - > Usage/API
- Scheduling
 - > Concepts
 - > Usage/API

(Italics describe things subject to change)



Threads

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Thread concept

Classic

- Unit of execution
- Lowest schedulable entity
- Preemptable
- Affine to a core
- Can block
- Can be synchronized with

Barrelfish

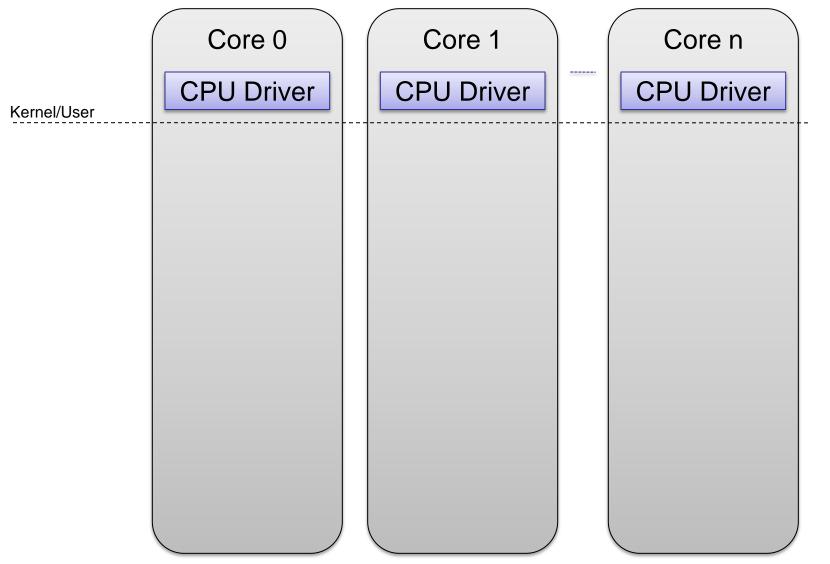
- Can invoke capabilities
- Can send/receive messages



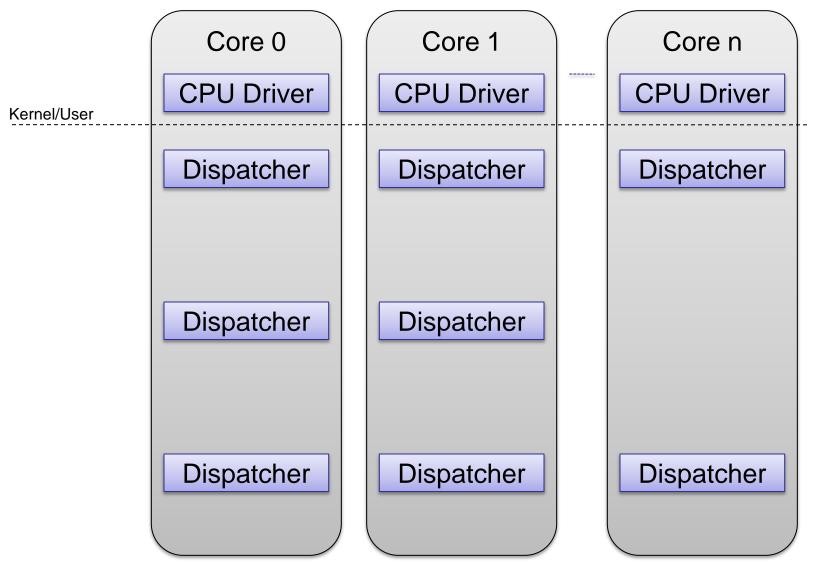
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Core 0	Core 1	Core n

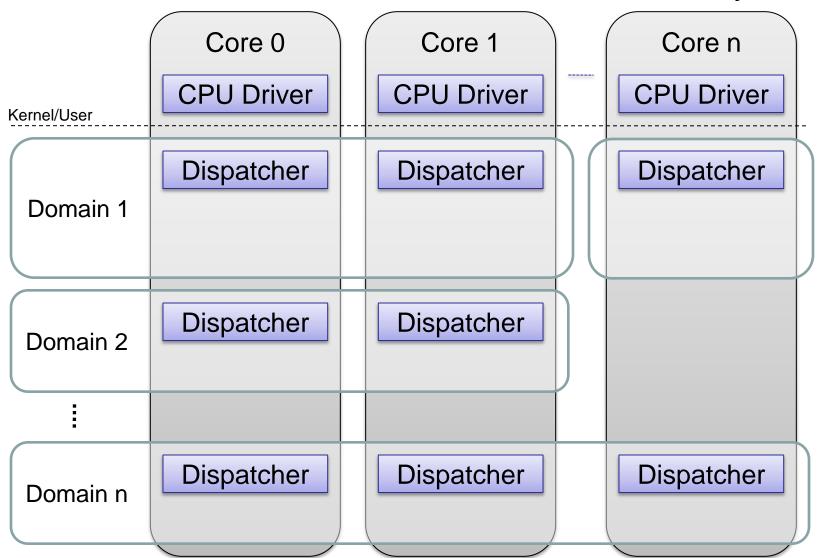




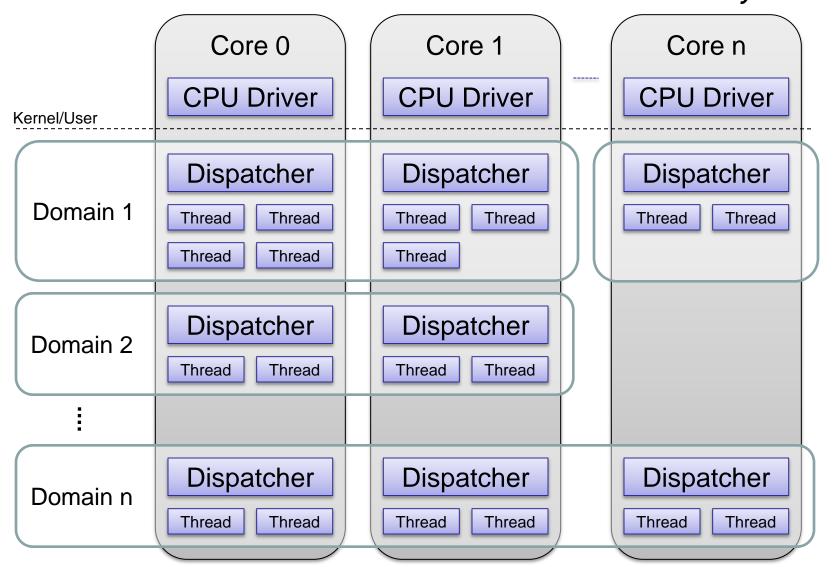




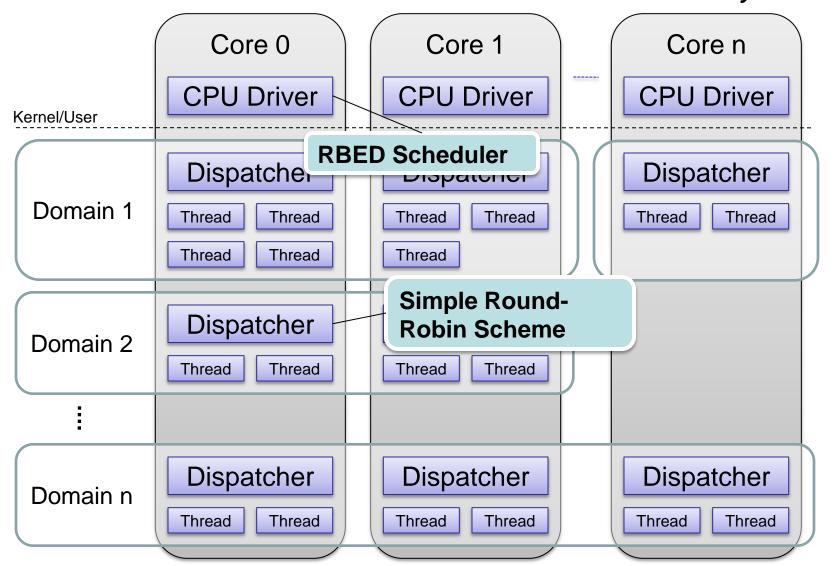














- Implemented in user-space
- Provided by a threads package
- Currently contained in libbarrelfish
 - > Public API in include/barrelfish/threads.h
 - Provides POSIX-like behavior

Thread data types



- struct thread *thread = thread_self();
 - > pthread_t equivalent, designates a thread
- **int** (*thread_func_t) (void *data)
 - Thread start routine
- struct thread mutex mutex = THREAD MUTEX INITIALIZER;
 - Mutex (can be nested and tested)
- struct thread_cond cond = THREAD_COND_INITIALIZER;
 - Condition variable (signal and broadcast semantics)
- struct thread_sem sem = THREAD_SEM_INITIALIZER;
 - Semaphore (can be tested)

Thread creation



• Create thread on local core

```
struct thread *thread_create(thread_func_t start_func,
void *data);
struct thread *thread_create_varstack(thread_func_t
```

start func, void *arg, size t stacksize);

Thread stack caveats

- Allocated at thread creation time
- Don't grow dynamically
- Not protected against overflow



- One void * can be associated with each thread
 - > Bad composability, we know
- Pthreads has thread local key-value store instead

void thread set tls(void *tls);

void *thread_get_tls(void);

Other useful thread operations



• Return thread ID

struct thread *thread_self(void);

Yield timeslice of this thread

void thread_yield(void);

Exit thread

void thread_exit(void);

- Threads are only cleaned up when joined with
 - Unless detached

errval_t thread_join(struct thread *thread, int *retval);
errval_t thread_detach(struct thread *thread);

What we don't have but POSIX does



- Forking-related operations (pthread_atfork())
- Attributes
 - Scheduling policy set via scheduler API
 - No guarded stacks
 - No contention scopes
- Barriers
 - Done elsewhere
- Cancelable threads
- Timed synchronization primitives
- Concurrency levels
- Thread signals
- Read/write locks

Multi-core Threading

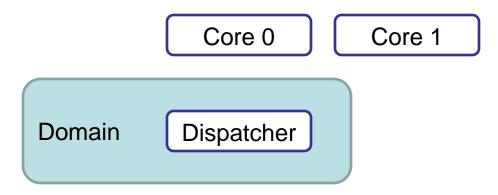


• "Span" your domain to other cores:

errval_t domain_new_dispatcher(uint8_t core_id,
domain_spanned_callback_t callback, void *callback_arg);

static void domain_spanned(void *arg, errval_t err);

• Asynchronous operation



Multi-core Threading

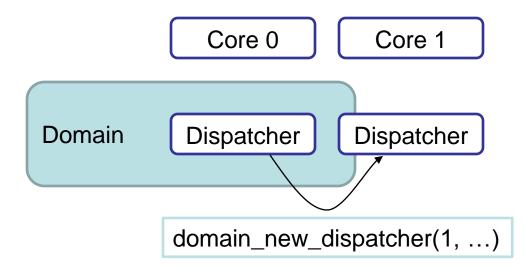


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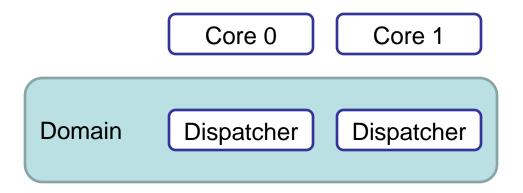


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Spanning a domain

- New dispatcher has its own resources
 - Monitor, memory server bindings
 - Slabs, slots and heaps
 - Page tables
 - Capabilities
 - All other bindings
- All other resources are shared
 - Virtual address space
 - Locks
 - Semaphores
 - Condition variables

▶ ...



Multi-core threading



• Start thread on different core:

```
errval_t domain_thread_create_on(coreid_t core_id,
thread_func_t start_func, void *arg);
```

- Move thread to different core:
 - Currently only supports self-migration
 - > Doesn't migrate open connections!

```
errval_t domain_thread_move_to(struct thread *thread,
coreid_t core_id);
```



Scheduling

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- Rate-Based Earliest Deadline First (RBED)
 - Algorithm by Scott Brandt, UC Santa Cruz
 - Deterministic, versatile, unified model
- Best-effort tasks with priorities
 - > UNIX-style I/O priority boost
- Rate-based tasks
 - > Worst-case execution time, deadline, period
- Soft-realtime
 - > Upcall when deadline missed
- Hard-realtime
 - > Admission control

Scheduling at multiple timescales



- Long-term
 - Scheduler manifests
 - Divide app into phases with RBED parameters
- Mid-term
 - Phase changes
- Short-term
 - Scheduled by RBED

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Scheduler API



- Public API in include/barrelfish/resource_ctrl.h
- Submit scheduling manifest on local dispatcher

```
errval_t rsrc_manifest(const char *manifest, rsrcid_t
*id);
```

• Join the manifest from other dispatchers

```
errval_t rsrc_join(rsrcid_t id);
```

• Change resource **phase** from any joined dispatcher

errval_t rsrc_phase(rsrcid_t id, uint32_t phase);



static const char *my_manifest =
 "B 1\n" // Normal phase
 "H 20 160 160\n"; // Hard real-time phase

- Phases are implicitly numbered from 0
- Phase 0: Best-effort with priority 1
- Phase 1: Hard real-time
 - > Worst-case execution time 20ms
 - Period 160ms
 - Deadline 160ms



Thank you

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Backup

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